

Problem 1. (DPV Exercise 7.1: 20 points)

See Figure 1. Let $obj(x,y) = 5x + 3y$. Then $obj(0,0) = 0$, $obj(5,0) = 25$, $obj(5,2) = 31$, and $obj(2,5) = 25$. So the max is 31 at $(5,2)$.

Problem 2. (DPV Exercise 7.4: 20 points)

Variables: x_r = pints of regular Duff ordered and x_s = pints of strong Duff ordered.

Objective function to maximize: $1 \cdot x_r + 1.5 \cdot x_s$.

Constraints: $2x_s \leq x_r$, $x_r + x_s \leq 3000$, $x_r, x_s \geq 0$.

See Figure 2. Let $obj(x_r, x_s) = x_r + 1.5 \cdot x_s$. Then: $obj(0,0) = 0$, $obj(3000,0) = 3000$, and $obj(2000,1000) = 3500$. So \$3,500 is the maximum at $(2000,1000)$.

Problem 3. (DPV Exercise 7.6: 8 points)

Minimize: $x + y$ under constraints $x, y \geq 0$. The feasible region is the entire 1st quadrant, but the $x + y$ has a unique minimal value at $(0,0)$.

Problem 4. (DPV Exercise 7.7: 12 points)

(a) In order to be infeasible, the half-plane described by $ax + by \leq 1$ does not intersect the 1st quadrant. But if $x = y = 0$, then $0 = ax + by \leq 1$. So there always a feasible point.

(b) $a \leq 0$ or $b \leq 0$.

(c) $a > 0$ and $b > 0$.

Problem 5. (DPV Exercise 7.17: 20 points)

(a) 11 is the max-flow.

$(\{S,A,B\}, \{C,D,T\})$ is the min-cut.

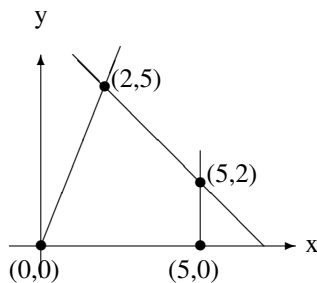


Figure 1

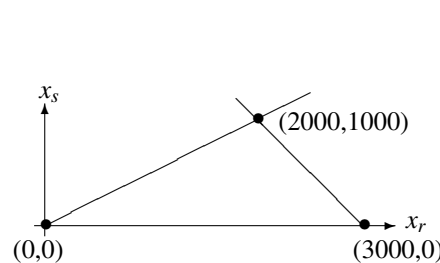


Figure 2

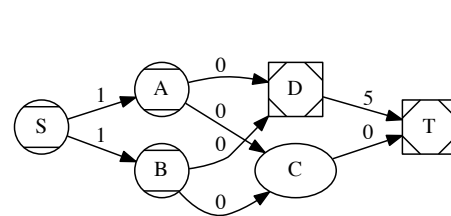


Figure 3

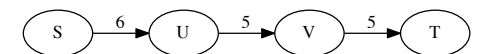


Figure 4

(b) See Figure 3. The nodes reachable from S are in funny circles and the nodes reachable from T are in funny squares.

(c) (A,C) and (B,C) .

(d) See Figure 4.

(e) Construct the residual graph and do a graph search that looks for edges with residual capacity 0 but have edges on both ends of them with positive residual capacity.

Problem 6. (DPV Exercise 7.18 parts (a) and (b): 20 points)

(a) From the original graph $G = (V, E)$, construct a graph $G' = (V', E')$ where:

$$V' = V \cup \{S_*, T_*\}.$$

$$E' = E \cup \{(S_*, s) : s \text{ is a source in } G\} \cup \{(t, T_*) : t \text{ is a target in } G\}$$

Each edge from E has the same capacity as in G . Each (S_*, s) and (t, T_*) edge has capacity C where C is, say, the sum of all the capacities in G . S_* and T_* are the source and target nodes in G' , respectively.

(b) From the original graph $G = (V, E)$, construct a graph $G' = (V', E')$ where:

$$V' = V \cup \{\hat{v} : v \in V\}.$$

$$E' = \{(u, \hat{v}) : (u, v) \in E\} \cup \{(\hat{v}, v) : v \in V\}.$$

Each (u, \hat{v}) has the same capacity as (u, v) in G . Each (\hat{v}, v) has the max-capacity of v . \hat{S} is the source node for G' .